Protein Design

- What is it?
- Why?
- Experimental methods
- What we need
- Computational Methods
- introduce
 - Monte Carlo
 - a pruning algorithm

What is protein design?

Assumption

- you can write a protein sequence on a piece of paper
- a molecular biologist can produce it

Most general

- you have a protein which is useful (enzyme, binding, ...)
- you want to make it more stable
 - temperature
 - solvents (tolerate organic solvents)
 - pH
 - we concentrate on stability

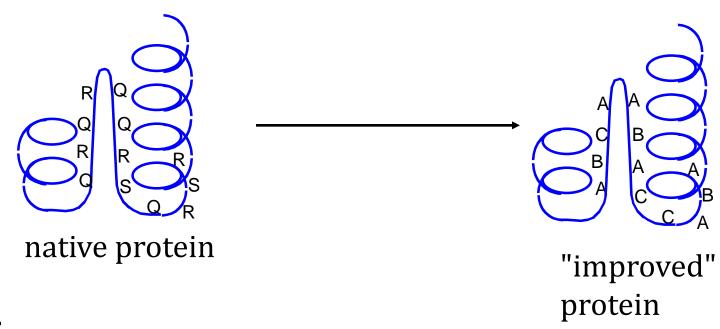
Experimental approaches

- Bacteria / selection
- For binding
 - phage display
 - in vitro evolution
- stability more difficult
- computational methods...

Formalising the problem

We have a working structure

want to make it more stable



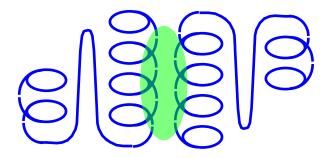
Rules

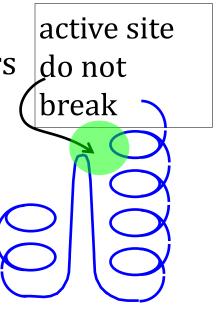
- structure should not change
- should be able to fix some residues (active site, important)...

Fixing / specifying residues

Examples

- lysine (K) often used for binding
 - change a residue to K and protein does not fold
 - mission:
 - adapt the rest of the residues to be stable
- change all residues, but not those in active site
- change some residues at surface to be soluble
- change some residues at surface to stop dimers





Ingredients

- Score function (like energy)
- Search method

Score function

- how does sequence fit to structure?
- sequence $S=\{s_1, s_2, ...s_N\}$
- coordinates $R = \{ \mathbf{r}_1, \mathbf{r}_2, \dots \mathbf{r}_N \}$
- score = f(S, R) (different nomenclature soon)
- mission
 - adjust S to as to maximise score (minimise quasi-energy)

Score function

How do amino acids

- suit structure? $score = \sum_{i=1}^{N_{res}} score_{struct}(s_i, R)$
- suit each other? $+ \sum_{i=1}^{N_{res}} \sum_{j>i}^{N_{res}} score_{pair}(s_i, s_j, R)$

score_{struct} might have

- backbone preferences (no proline in helices, ..)
- solvation (penalise hydrophobic at surface)

$score_{pair}$

- are residues too big (clashing)
- are there holes? charges near each other?

Messy functions

lots of parameters

Searching

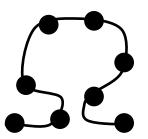
Systematic search – how long?



• search space for $N_{res} = 20 \times 20 \times \dots = 20^{N_{res}}$

Search space complex

- every time you change a residue, affects all neighbours
 - effects neighbours of neighbours
- brute force not a good idea
- two methods here
 - 1. Monte Carlo / simulated annealing
 - 2. Pruning / dead end elimination



Monte Carlo

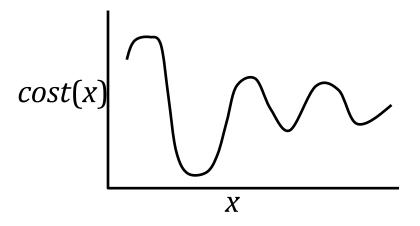
- more formally next semester
- first the problem

The sequence optimisation problem

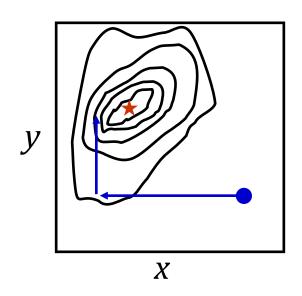
- discrete
- local minima / correlations in surface
- high dimensional

dimensions and correlations

• a 1D problem



- local minima
- minimum of *x* depends on *y*
- cannot optimize x and y independently
- what are correlations in this problem?



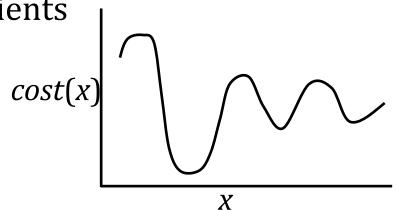
Discrete vs continuous problems

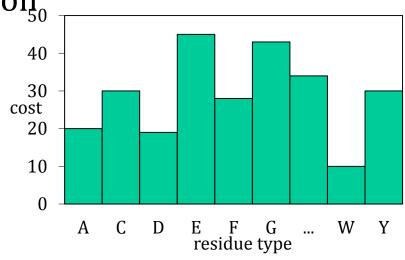
For a continuous function use gradients

- to optimise
- to recognise minima / maxima
- continuous functions
 - step in one direction is good
 - try another in same direction



- no gradients
- order of labels arbitrary
 - ACDE or ECAD
- discrete
 - step in one direction may be no predictor of best direction





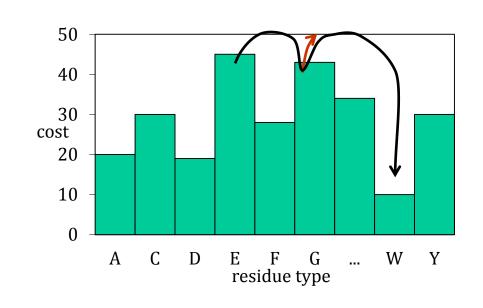
what do we want?

From step to step (sequence to sequence)

- be prepared to move in any direction
- if the system improves, try not to throw away good properties
- must be willing to go uphill sometimes

Philosophy

- take a random move
- if it improves system
 - keep it
- if cost becomes worse
 - sometimes keep it
 - sometime reject



Acceptance / rejection

- for convenience, write $cost(S_n)$ neglect coordinates R Sign convention
- system (sequence) at step n is S_n
- after a random step, cost changes from $cost(S_n)$ to $cost(S_{n+1})$
- $\Delta c = cost(S_{n+1}) cost(S_n)$
- our sign convention: if Δc < 0, system is better

When to accept?

- if Δc is a bit < 0 accept
- if Δc is a bit > 0, maybe OK
- if $\Delta c >> 0$, do not accept

Formal acceptance rule

- $-\Delta c < 0$, $e^{-\Delta c}$ is between 0..1
- $-\Delta c \approx 0$ then $e^{-\Delta c} \approx 1$ as $\Delta c \rightarrow \infty$ then $e^{-\Delta c} \rightarrow 0$
- formalise this rule

```
set up S=S_0 and cost(S_0)

while (not finished)

S_{trial} = random \ step \ from \ S

\Delta c = cost(S_{trial}) - cost(S)

if (\Delta c < 0) /* accept */

S=S_{trial}

else

r = rand \ (0..1)

if (e^{-\Delta c} \ge r)

S=S_{trial}
```

vorsicht! not the final method

why we need temperature

As described

- system will run around
- try lots of new configurations
- sometimes accept bad moves
- always take good moves
- may never find best solution
 - imagine you are at a favourable state
 - most changes are uphill (unfavourable)
 - many of the smaller ones will be accepted
 - if we were to find the best sequence, the system would move away from it
- how to fix ?

why we need temperature

- Initial sequence is not so good
 - let the system change a lot and explore new possibilities
- after some searching, make the system less likely to go uphill
- introduce the concept of temperature *T*
- initially high T means you can go uphill (like a high energy state)
- as you cool the system down, it tends to find lowest energy state
- change acceptance criterion to $e^{\frac{-\Delta c}{T}}$ as

$$T \to \infty$$
, $e^{\frac{-\Delta c}{T}} \to 1$
 $T \to 0$, $e^{\frac{-\Delta c}{T}} \to 0$

• put this into previous description

why we need temperature

```
set up S = S_0 and cost(S_0) set T = T_0
while (not finished)
     S_{trial} = random step from S
                                            /* ε bit smaller than 1 */
     T = \epsilon T
     \Delta c = cost(S_{trial}) - cost(S)
     if (\Delta c < 0)
             S= S<sub>trial</sub>
     else
             r = rand (0..1)
             if (\exp(-\Delta c/T) \ge r)
                      S= S<sub>trial</sub>
```

Name of this procedure

"simulated annealing"

Final Monte Carlo / annealing

History applications

- discrete problems travelling salesman, circuit layout
- deterministic? No
- convergence ? Unknown

Practical issues

- what is a random step?
 - change one amino acid? change interacting pairs?
- easy to program
- lots of trial and error
- statistical properties next semester
- can we reduce the search space?

Pruning

Are there elements of sequence which are impossible?

at position 35, no chance of Y, W, I, L, ...

Can one find impossible combinations

 reduce the search space so it can be searched systematically (brute force)

... dead end elimination method

• use an energy-like nomenclature

Nomenclature

- we are not dealing with
 - free energy G or F or potential energy U or E
- but let us pretend
 - score is *E*
- rule : more negative *E* , better the system
- structure is fixed so neglect R / r terms
- define a function $s_i(a)$ as the residue type at site i
 - can take on 20 values of "a" why? foreach (a in A, C, D, E.., W, Y) evaluate energy corresponding to a
- our energies ?
 - two parts pairwise and residue with backbone

Nomenclature

E is (quasi-energy) of whole system

- label E_1 as the terms that depend on residue + fixed environment
- E_2 as the energy terms that depend on pairs

$$E = \sum_{i=1}^{N_{res}} E_1(s_i) + \sum_{i=1}^{N_{res}} \sum_{j \neq i}^{N_{res}} E_2(s_i, s_j)$$

If we are interested in site *i* and being in state *a* what do we have to look at ?

$$\sum_{i=1}^{N_{res}} E_1(s_i(a)) + \sum_{i=1}^{N_{res}} \sum_{j \neq i}^{N_{res}} E_2(s_i(a), s_j(b))$$

Nomenclature and rules

- there are $20 (N_{type})$ residues
- $\min E_1(s_i(a))$ which fits best to the fixed environment?
- implies testing each of the N_{type} for a
- what is the best energy type a at site i could have, interacting with one site *j*?

$$E_1(s_i(a)) + \min_b E_2(s_i(a), s_j(b))$$

 what is the best energy that type a at i could have considering all neighbours?

$$E_1(s_i(a)) + \sum_{j \neq i} \min_b E_2(s_i(a), s_j(b))$$

- $E_1(s_i(a)) + \sum_{j \neq i} \min_b E_2(s_i(a), s_j(b))$ for each a can work out what is the best score it could yield
 - loop over *b*
 - within loop over *j*

Dead-end elimination method

- worst energy that type c at i could have considering all neighbours? $E_1(s_i(c)) + \sum_{i \neq i} \max_d E_2(s_i(c), s_j(d))$
- when can one eliminate (rule out) residue type *a* at site *i*?
- for any residues a, c
- if the best energy for *a* is worse than the worst for *c*
 - a cannot be part of the optimal solution ... if

$$E_1(s_i(a)) + \sum_{j \neq i} \min_b E_2(s_i(a), s_j(b)) > E_1(s_i(c)) + \sum_{j \neq i} \max_d E_2(s_i(c), s_j(d))$$

Dead-end elimination method

$$\begin{split} E_1(s_i(a)) + \sum_{j \neq i} \min_b E_2(s_i(a), s_j(b)) > E_1(s_i(c)) + \sum_{j \neq i} \max_d E_2(s_i(c), s_j(d)) \\ \text{using this approach} \\ \text{for } (i = 0; i < N_{\text{res}} ; i++) \\ \text{foreach } a \text{ in } N_{\text{type}} \\ \text{calculate worst score for } a \\ \text{calculate best score for } a \\ \text{for } (i = 0; i < N_{\text{res}} ; i++) \\ \text{foreach } a \text{ in } N_{\text{type}} \\ \text{foreach } b \text{ in } N_{\text{type}} \\ \text{if best}(a) > \text{worst } (b) \\ \text{remove } a \text{ from candidates} \end{split}$$

How strong is this condition?

DEE condition

- much of the time
 - cannot really rule out type a
- example?
 - initial
 - 2×10^{27}
 - final
 - searchable in 90 cpu hr

Dahiyat, B.I, Mayo, S.L. (1997), Science 278, 82-87

deterministic

Combining ideas

- use DEE to get a list of candidate residues at each position
- search remaining space with Monte Carlo / simulated annealing
- not deterministic

Success

Method

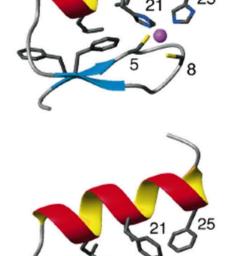
• Dead end elimination + systematic search

designed QQYTAKIKGRTFRNEKELRDFIEKFKGR

native KPFQCRICMRNFSRSDHLTTHIRTHTGE

New sequence

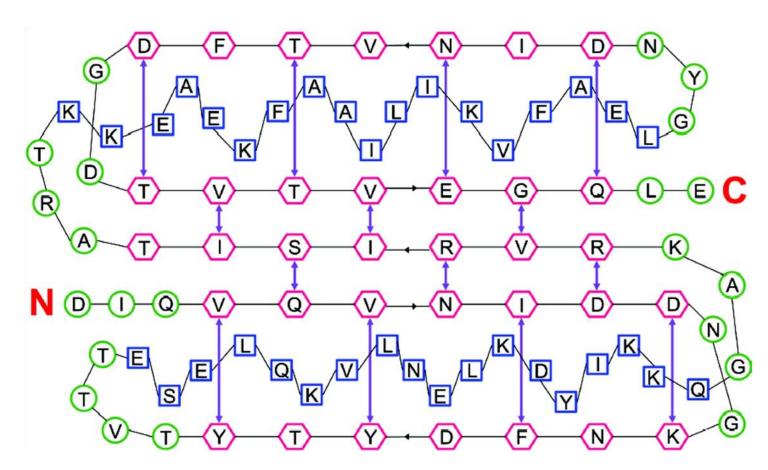
- about 20 % similar to start
- not related to any known protein (still)
- Structure solved by NMR
- Problem solved?
 - maybe not



Success

Mission

- sketch a new protein topology
- build a sequence to fit it



Success

Methods

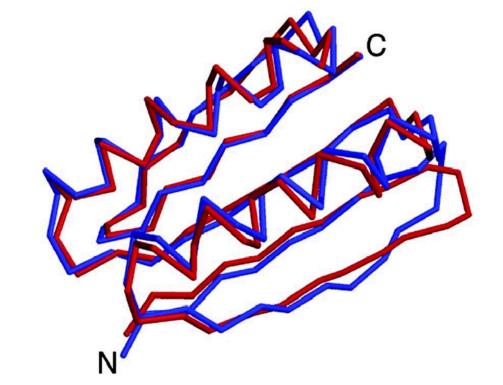
pure Monte Carlo

Result

apparently new sequence

Structure

- as predicted
- solved by X-ray
 - phasing story



- Problem solved
 - unclear (how many failures ?)

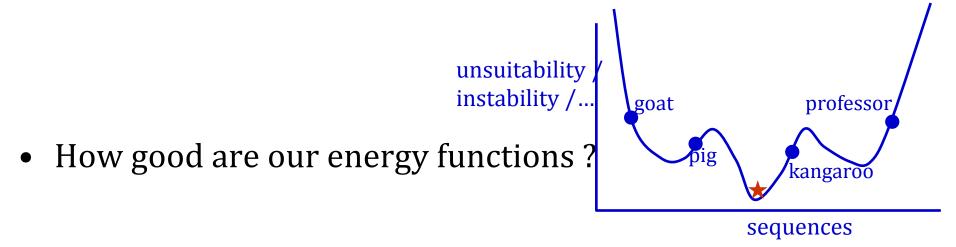
Methods so far

	Monte Carlo	Dead-end elimination
guaranteed global optimum	no	does not try
deterministic	no	yes

Only one answer?

May not matter

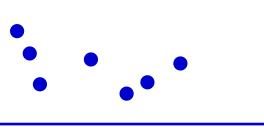
- consider real proteins compare human, goat, ...
 - all stable all slightly different
- implication
 - there may be many solutions which are equally good



Determinism and energy

unsuitability , instability /...

I have a perfect score / energy function

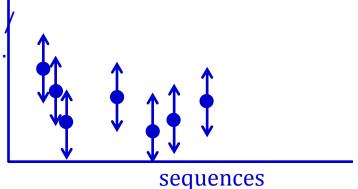


sequences

unsuitability instability /..

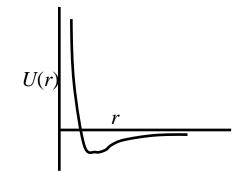
I have errors / approximations

best answer could be any one



Problems – stability / energy

- energy functions
- what do we mean by energy?
- example two charges $U(r) = \frac{q_1 q_2}{Dr}$
- example two argon atoms $U(r) = 4\varepsilon(\sigma^{12}r^{-12} \sigma^6r^{-6})$
- make energy better?
 - replace every amino acid by a larger one (more contacts – more negative energy)



- silly proteins are not full of large amino acids
- what determines stability?

Problems – stability / energy

- stability does a molecule prefer to be folded or unfolded ?
- what is unfolded? or ?

My energy function tells me to change "X" to "Y"

- it affects both the good 🔞 and bad 🔨
- has it affected the energy difference?
 - no guarantee

Current score functions?

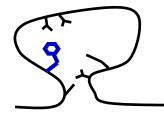
- some pure potential energy
- very difficult to estimate ΔG

Problems - sidechains

Side chain positions

- can I ever calculate the energy if I change X to Y?
- insert a phe into this structure
- what interactions does it have?





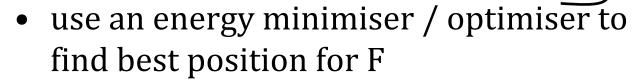


How to cope with side chain positions in a practical way

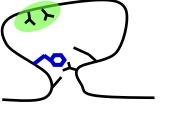
- optimise location of sidechains
- use average
- explicit rotamers

Sidechains - optimise at each step

- I start with known protein
 - change $A \rightarrow F$



- sensible?
 - we have a gigantic search space
 - explicit optimisation of one side chain would be expensive
- silly
- I change $A \rightarrow F$, but the rest of the side chains may move
- bad idea



Sidechains - use averaging

Ignore the problem of sidechain geometry

- at room temperature, side chains move
 - small (middle of protein) to big (surface)
- we cannot expect Å accuracy anyway
- rather fast
- what if we want to worry about atoms?

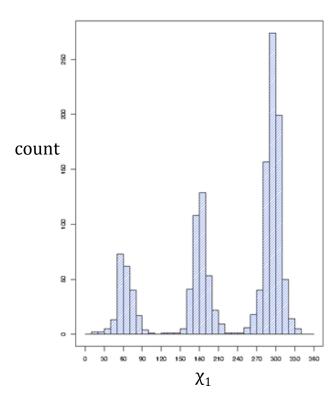
Sidechains - use rotamers

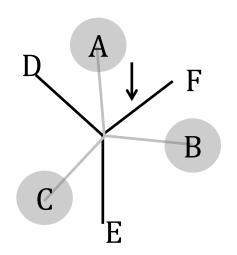
• sidechains can move anywhere but

• there are preferences in diagram – three more likely states

• how many times is the first angle (χ_1) seen at each angle ?

- how to use this?
 - look for most popular angles (60, 180, 300)



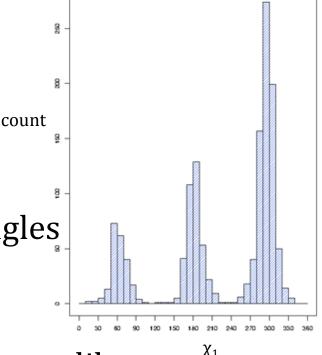


Sidechains - use rotamers

- For this example
 - do not have 1 cys residue
 - replace with cys1, cys2, cys3
 - treat all amino acids similarly
- more complicated because of more angles



- N_{type} of amino acids >> 20
- requires that you have a pre-built rotamer library



- fits to
 - Monte Carlo (random moves between residues or rotamers)
 - dead end elimination (will remove impossible rotamers)

Problems - viability

Designed sequences must

- fold
- be expressed + produced

Summary

- Nature of the problem discrete (not continuous)
- Optimisation methods (MC, DEE)
- Score functions
 - not energy, not free energy, not potential energy
- Success / state of the art
 - not many examples from literature
 - failure rate?
 - cost
- Definitely not a routine method